Using Natural Deduction to Represent Arguments (Part 2)

Scott Martin

February 15, 2012

More Arguments and Rules

Example Argument

Consider the following:

(1)
a. Pastor Ingqvist and Father Wilmer go fishing.
b. If Pastor Ingqvist goes Fishing, no one receives the lutefisk shipment.
c. If no one receives the lutefisk shipment and today is Saturday, the festival is canceled.
d. Today is Saturday.
e. That means the festival must be canceled.

Analyzing the Example

To analyze the argument in (1) for validity, we proceed as usual, starting by identifying the atomic propositions it uses.

P Pastor Ingqvist goes fishing.

W Father Wilmer goes fishing.

L Someone receives the lutefisk shipment.

S Today is Saturday.

C The festival is canceled.

Translating the Argument into PL

Now we can translate the declarative sentences used to construct the argument in (1) into propositions:

$$P \wedge W \tag{1a}$$
$$P \rightarrow \neg L \tag{1b}$$
$$(\neg L \wedge S) \rightarrow C \tag{1c}$$
$$S \tag{1d}$$
$$C \tag{1e}$$

So as usual, we have an argument with some premises (1a-1d) and a conclusion (1e).

Strategy for Giving a Formal Proof

- However, notice that the instances of $\wedge$ complicate things somewhat.
- In order to make the inference step that lets us use $P \to \neg L$ to get $\neg L$, we need a proof of the antecedent P.
- But our assumptions only have a proof of $P \wedge W$.
- Similarly, we need to prove $\neg L \wedge S$ in order to conclude C.
- But the inference step we'll use to go from $P \to \neg L$ to $\neg L$ given P will only give us a proof of $\neg L$ by itself.
- We know, both intuitively and via truth table verification that A being true and B being true means that $A \wedge B$ is also true.
- And while we can sometimes just add more assumptions containing just the left and right sides of a conjunction, this won't work all the time.
- We need more rules to handle this argument using natural deduction.

Rule for Introducing $\wedge$

As before, φ and ψ are meta-variables ranging over propositions (atomic or complex).

Inference Rule 4 (Conjunction Introduction)**.**

$$\frac{\varphi \qquad \psi}{\varphi \wedge \psi}\ (\wedge\mathrm{I})$$

- Rule 4 is called an introduction rule because it introduces an instance of the connective $\wedge$ where one was not present before.
- It says that if you've proved φ and you've proved ψ, then you've proved $\varphi \wedge \psi$.
- Rule 4 is both intuitively straightforward and easily verifiable using a truth table.

Rules for Eliminating $\wedge$

To "unpack" a conjunction into its component parts, we need two rules that essentially do the same thing:

Inference Rule 5 (Conjunction Elimination 1)**.**

$$\frac{\varphi \wedge \psi}{\varphi}\ (\wedge\mathrm{E}_{1})$$

Inference Rule 6 (Conjunction Elimination 2)**.**

$$\frac{\varphi \wedge \psi}{\psi}\ (\wedge\mathrm{E}_{2})$$

- Rules 5 and 6 are mirror images of each other.
- They say that if you've proved the conjunction $\varphi \wedge \psi$ then you can deduce that you've proved either of the conjuncts.

$$\dfrac{\dfrac{\dfrac{\dfrac{\dfrac{}{P \wedge W}\,(\text{Hyp})}{P}\,(\wedge\text{E}_1) \qquad \dfrac{}{P \to \neg L}\,(\text{Hyp})}{\neg L}\,(\to\text{E}) \qquad \dfrac{}{S}\,(\text{Hyp})}{\neg L \wedge S}\,(\wedge\text{I}) \qquad \dfrac{}{(\neg L \wedge S) \to C}\,(\text{Hyp})}{C}\,(\to\text{E})$$

Figure 1: Proof of the argument in (1).

Applying the new rules to the argument

- With Rules 4-6, we have everything we need to both combine proofs via conjunction and separate conjoined parts into their two pieces.
- So, given that we already have a way to eliminate the $\to$ connective, Figure 1 contains a formal proof of the argument in (1).

Homework

Exercises

Problem 1. We know, both intuitively and from truth tables, that for any two propositions φ and ψ the propositions $\varphi \wedge \psi$ and $\psi \wedge \varphi$ are equivalent. Give a formal proof that has $A \wedge B$ as its premise and $B \wedge A$ as its conclusion. That is, you should come up with a proof tree that looks like

$$\dfrac{}{A \wedge B}\,(\text{Hyp})$$

$$\dfrac{\vdots}{B \wedge A}\,(?)$$

where you fill in the $\vdots$ and ?s. (Hint: you will use the rules for $\wedge$ talked about above.)